

Goals You learn to formalize simple situations and recursive functions in Isabelle and prove simple theorems about these.

Hint On this sheet, all proofs can be found by `auto`, if you use induction where necessary and give the right equations to the simplifier.

Exercise 1 [4] Functional Programming

Write a function

```
sumup :: "int list ⇒ int"
```

which computes the sum of numbers of a given Isabelle/HOL list.

Evaluate the function on concrete values to get a feeling for its correctness:

```
value "sumup []"  
value "sumup [1]"  
value "sumup [1,2]"  
value "sumup [1,2,45,62]"
```

Prove now that the `sumup` function is distributive, i.e. that `sumup (xs @ ys) = sumup xs + sumup ys` holds.

Afterwards, prove that `sumup` returns the same result for reversed lists, i.e. that `sumup (rev xs) = sumup xs` holds

Exercise 2 [2] Nodes of binary trees

In the exercise theory, we defined a type of binary trees with functions `nleaves` and `ninner` counting the leaves and inner nodes.

Define a predicate

```
fun bt-full :: "bt ⇒ bool"
```

recognizing full binary trees. A binary tree is called *full*, iff for all Nodes either both or none of the children are Tip.

Formulate and prove the property that for each full tree holds `leaves t = ninner t + 1`.

Hint: To prove this property, you will need the following lemma:

```
t ≠ Tip ↔ (∃ l v r. t = Node l v r)
```

This can be proved by case analysis on `t`, which can be performed with `apply (cases t)`.

Exercise 3 [4] Summing Binary Trees

Write a function

```
sumtree :: "bt ⇒ int"
```

which computes the sum of all elements of a tree and a function

```
list-of-tree :: "bt ⇒ int list"
```

which computes the list of all elements (including duplicates) of a tree.

Then prove `sumtree xs = sumup (list-of-tree xs)`.
