

### Exercise 1 (Normal Forms)

Recall the inductive definition of the set NF of *normal forms*:

$$\frac{t \in \text{NF}}{\lambda x. t \in \text{NF}}$$

$$\frac{n \geq 0 \quad t_1 \in \text{NF} \quad t_2 \in \text{NF} \quad \dots \quad t_n \in \text{NF}}{x t_1 t_2 \dots t_n \in \text{NF}}$$

Show that this set precisely captures all normal forms, i.e.:

$$t \in \text{NF} \Leftrightarrow \nexists t'. t \rightarrow_\beta t'$$

### Solution

**Direction  $\Rightarrow$ :** Proof by induction on the derivation of  $t \in \text{NF}$ .

**Case  $t = \lambda x. s \in \text{NF}$**

We obtain the induction hypothesis  $s \in \text{NF} \Rightarrow \nexists s'. s \rightarrow_\beta s'$ . Suppose there exists an  $s'$  such that  $\lambda x. s \rightarrow_\beta \lambda x. s'$ . By rule inversion, it must hold that  $s \rightarrow_\beta s'$  which is a contradiction to induction hypothesis where we discharge the assumption  $s \in \text{NF}$  by rule inversion on  $t \in \text{NF}$ .

**Case  $t = x t_1 \dots t_n \in \text{NF}$**

By rule inversion on  $t \in \text{NF}$  we get  $t_i \in \text{NF}$  and therefore  $\nexists t'. t_i \rightarrow_\beta t'$  by the IH for  $1 \leq i \leq n$ . We show this case by another induction on  $n$ . The case  $n = 0$  is trivial. Now assume  $\nexists t'. x t_1 t_2 \dots t_n \rightarrow_\beta t'$  as the induction hypothesis, and  $x t_1 t_2 \dots t_n t_{n+1} \rightarrow_\beta t'$  for the sake of contradiction. By analysing the derivation of the latter, we can only conclude  $\exists t'. t_{n+1} \rightarrow_\beta t'$ , which manifests a contradiction.

**Direction  $\Leftarrow$ :** We show  $t \notin \text{NF} \Rightarrow \exists t'. t \rightarrow_\beta t'$  by structural induction on  $t$ .

**Case  $t = x$**

Then  $t \in \text{NF}$  which is a contradiction to the assumption  $t \notin \text{NF}$ .

**Case  $t = \lambda x. s$**

We assume that  $t = \lambda x. s \notin \text{NF}$ . Thus, it must hold that  $s \notin \text{NF}$  by rule inversion, which together with the IH lets us obtain a  $s'$  where  $s \rightarrow_\beta s'$ . Hence, we conclude that  $\lambda x. s \rightarrow_\beta \lambda x. s'$ .

**Case**  $t = t_1 t_2$

We assume  $t_1 \notin \text{NF} \implies \exists t'. t_1 \rightarrow_\beta t'$  and  $t_2 \notin \text{NF} \implies \exists t'. t_2 \rightarrow_\beta t'$  as the induction hypothesis, and  $t \notin \text{NF}$ .

The cases where  $t_1 \notin \text{NF}$  or  $t_2 \notin \text{NF}$  are immediate by the induction hypotheses together with the definition of  $\rightarrow_\beta$ .

Consider the case  $t_1, t_2 \in \text{NF}$ . We analyze the derivation of  $t_1 \in \text{NF}$ .

**Case**  $t_1 = x s_1 \dots s_n$  **where**  $s_1, \dots, s_n \in \text{NF}$

It follows that  $t = t_1 t_2 = x s_1 \dots s_n t_2 \in \text{NF}$  since  $t_2 \in \text{NF}$ . But this is a contradiction to the assumption  $t_1 t_2 \notin \text{NF}$ .

**Case**  $t_1 = (\lambda x. s)$  **where**  $s \in \text{NF}$

Then it holds that  $t = t_1 t_2 = (\lambda x. s) t_2 \rightarrow_\beta s[t_2/x]$ .

## Exercise 2 (Weakly Normalising Terms)

Inductively define the set of weakly normalising terms WN, i.e. the set of terms that have a  $\beta$ -normal form. In particular it should hold that

$$s \in \text{WN} \iff \exists t. s \Rightarrow_n t.$$

Similarly, define the set of strongly normalising terms SN where a term  $s$  is strongly normalising if there is no infinite sequence  $\{t_i \mid i \in \mathbb{N}\}$  with  $s \rightarrow_\beta^* t_0$  and  $t_i \rightarrow_\beta t_{i+1}$  for  $i \in \mathbb{N}$ .

Give a term  $t$  that is weakly but not strongly normalising.

## Solution

We define WN as follows:

1.  $t_1, \dots, t_m \in \text{WN} \implies x t_1 \dots t_m \in \text{WN}$
2.  $t \in \text{WN} \implies \lambda x. t \in \text{WN}$
3.  $s[t/x] t_1 \dots t_m \in \text{WN} \implies (\lambda x. s) t t_1 \dots t_m \in \text{WN}$

The definition of SN is identical with exception of the last rule where we also demand that  $t \in \text{SN}$ .

Let  $\omega := (\lambda x. x x)$ . Then, it holds that  $(\lambda x. y) (\omega \omega) \in \text{WN} \setminus \text{SN}$ .

### Homework 3 (Normal Forms & Big Step)

Show:

$$t \in \text{NF} \wedge t \Rightarrow_n u \implies u = t$$

### Homework 4 (Characterisation of WN)

Prove the following characterisation of WN:

$$t \in \text{WN} \iff \exists t'. t \rightarrow_{\beta}^* t' \wedge t' \in \text{NF}$$

You may use the fact that

$$s \in \text{WN} \iff \exists t. s \Rightarrow_n t$$

and all theorems up to Theorem 1.5.8 from the lecture.