

LOGIC EXERCISES

TECHNICAL UNIVERSITY OF MUNICH
CHAIR FOR LOGIC AND VERIFICATION

PROF. TOBIAS NIPKOW
KEVIN KAPPELMANN

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EXERCISE SHEET 12

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The tutorial takes place on 06.07, 12–14.

Exercise 12.1. [Łoś–Vaught Test]

Let T be an L -theory with no finite models. Let $\kappa \geq |L|$ be a cardinal. Show that if any two models of size κ for T are elementarily equivalent, then T is complete.

You can assume the following without a proof:

Theorem 1 (Generalised Löwenheim-Skolem Theorems). *Let S be a set of formulas in a language of cardinality λ , and assume that S has some infinite model. Then for every infinite cardinal $\kappa \geq \lambda$, there is a model of cardinality κ for S .*

Exercise 12.2. [QE for DLO]

Use the quantifier-elimination procedure for DLOs to check whether the following formula is a member of $Th(DLO)$:

$$\exists x \forall y \exists z ((x < y \vee z < x) \wedge y < z)$$

Exercise 12.3. [Fourier–Motzkin Elimination]

Apply the Fourier–Motzkin Elimination to check the following sentences:

1. $\exists x \exists y (2 \cdot x + 3 \cdot y = 7 \wedge x < y \wedge 0 < x)$
2. $\exists x \exists y (3 \cdot x + 3 \cdot y < 8 \wedge 8 < 3 \cdot x + 2 \cdot y)$

Exercise 12.4. [Ferrante–Rackoff Elimination]

Apply the Ferrante–Rackoff Elimination to check the following sentence:

$$\exists x (\exists y (x = 2 \cdot y) \rightarrow (2 \cdot x \geq 0 \vee 3 \cdot x < 2))$$

Homework 12.1. [Subtraction Logic] (+++)

We consider a fragment of linear arithmetic, in which atomic formulas only take the form $x - y \leq c$ for variables x and y , and $c \in \mathbb{R}$.

For a finite set S of such difference constraints, we can define a corresponding inequality graph $G(V, E)$, where V is the set of variables of S , and E consists of all the edges (x, y) with weight c for all constraints $x - y \leq c$ of S . Show that the conjunction of all constraints from S is satisfiable iff G does not contain a negative cycle.

How can you use this theorem to obtain a procedure for deciding whether a formula is a member of this fragment where all variables and constants are of the domain \mathbb{Z} ?

Homework 12.2. [Min, Max, Abs] (++)

1. Show that $\text{Th}(\mathbb{R}, 0, 1, <, =, +, \min, \max)$ is decidable, where \min and \max return the minimum and maximum of two values.
2. Show that $\text{Th}(\mathbb{R}, 0, 1, <, =, +, \min, \max, |\cdot|)$ is decidable, where $|\cdot|$ is the absolute value.

Homework 12.3. [Optimising DLO] (++)

DLO suffers from a heavy performance loss because after each step, a DNF needs to be reconstructed. We want to study an optimisation that may avoid this under some circumstances.

Assume that we want to eliminate an $\exists xF$ where

- F contains no negations and quantifiers, and
- F contains no \perp ,
- there are only lower or only upper bounds for x in F .

Then, $\exists xF \equiv \top$. Prove the correctness of this optimisation.

In order to attain the impossible, one must attempt the absurd.

— Miguel de Cervantes