

Semantics of Programming Languages

Exercise Sheet 13

Exercise 13.1 Inverse Analysis

Consider a simple sign analysis based on this abstract domain:

```
datatype sign = None | Neg | Pos0 | Any
```

```
fun  $\gamma$  :: "sign  $\Rightarrow$  val set" where  
"  $\gamma$  None = {} " |  
"  $\gamma$  Neg = {i. i < 0} " |  
"  $\gamma$  Pos0 = {i. i  $\geq$  0} " |  
"  $\gamma$  Any = UNIV "
```

Define inverse analyses for "+" and "<" and prove the required correctness properties:

```
fun inv_plus' :: "sign  $\Rightarrow$  sign  $\Rightarrow$  sign  $\Rightarrow$  sign * sign"
```

lemma

```
"  $\llbracket$  inv_plus' a a1 a2 = (a1', a2'); i1  $\in$   $\gamma$  a1; i2  $\in$   $\gamma$  a2; i1+i2  $\in$   $\gamma$  a  $\rrbracket$   
 $\implies$  i1  $\in$   $\gamma$  a1'  $\wedge$  i2  $\in$   $\gamma$  a2' "
```

```
fun inv_less' :: "bool  $\Rightarrow$  sign  $\Rightarrow$  sign  $\Rightarrow$  sign * sign"
```

lemma

```
"  $\llbracket$  inv_less' bv a1 a2 = (a1', a2'); i1  $\in$   $\gamma$  a1; i2  $\in$   $\gamma$  a2; (i1 < i2) = bv  $\rrbracket$   
 $\implies$  i1  $\in$   $\gamma$  a1'  $\wedge$  i2  $\in$   $\gamma$  a2' "
```

Homework 13.1 Inverse Analysis for the Extended Reals

Submission until Monday, Feb 6, 23:59pm.

Extend the abstract interpreter from the last sheet with an inverse analysis. We skip the technical part, so you only need to define the inverse operations:

```
fun inv_plus' :: "bounds  $\Rightarrow$  bounds  $\Rightarrow$  bounds  $\Rightarrow$  (bounds * bounds)"
```

```
fun inv_less' :: "bool option  $\Rightarrow$  bounds  $\Rightarrow$  bounds  $\Rightarrow$  (bounds * bounds)"
```

They should be as precise as possible, i.e., the result minimal. For example:

```
value "inv_plus' (B {NaN}) (B { $\infty^+$ }) (B { $\infty^-$ , Real}) = (B { $\infty^+$ }, B { $\infty^-$ })"
```

```
value "inv_plus' (B { $\infty^+$ ,  $\infty^-$ }) (B { $\infty^+$ , NaN}) (B { $\infty^-$ , Real,  $\infty^+$ }) = (B { $\infty^+$ }, B { $\infty^+$ , Real})"
```

```
value "inv_less' (Some True) (B { $\infty^+$ ,  $\infty^-$ }) (B { $\infty^+$ }) = (B { $\infty^-$ }, B { $\infty^+$ })"
```

Prove them sound!

lemma *inv_plus'_sound*:

assumes “*inv_plus'* (B A) (B A1) (B A2) = (B A1', B A2')”
and “*i1* ∈ γ_bounds (B A1)”
and “*i2* ∈ γ_bounds (B A2)”
and “(case (*i1*,*i2*) of
 (*None*, *_*) ⇒ *None* ∈ γ_bounds (B A)
 (*_*, *None*) ⇒ *None* ∈ γ_bounds (B A)
 (*Some PInfty*, *Some MInfty*) ⇒ *None* ∈ γ_bounds (B A)
 (*Some MInfty*, *Some PInfty*) ⇒ *None* ∈ γ_bounds (B A)
 (*Some i1*, *Some i2*) ⇒ *Some* (*i1+i2*) ∈ γ_bounds (B A))”
shows “*i1* ∈ γ_bounds (B A1') ∧ *i2* ∈ γ_bounds (B A2')”

lemma *inv_less'_sound*:

assumes “*inv_less' bv* (B A1) (B A2) = (B A1', B A2')”
and “*i1* ∈ γ_bounds (B A1)”
and “*i2* ∈ γ_bounds (B A2)”
and “(case (*i1*,*i2*) of
 (*None*, *_*) ⇒ *None* = *bv*
 (*_*, *None*) ⇒ *None* = *bv*
 (*Some i1*, *Some i2*) ⇒ *Some* (*i1<i2*) = *bv*)”
shows “*i1* ∈ γ_bounds (B A1') ∧ *i2* ∈ γ_bounds (B A2')”